

University of Mumbai

Examinations Commencing from 7th January 2021 to 20th January 2021

Program: Computer Engineering

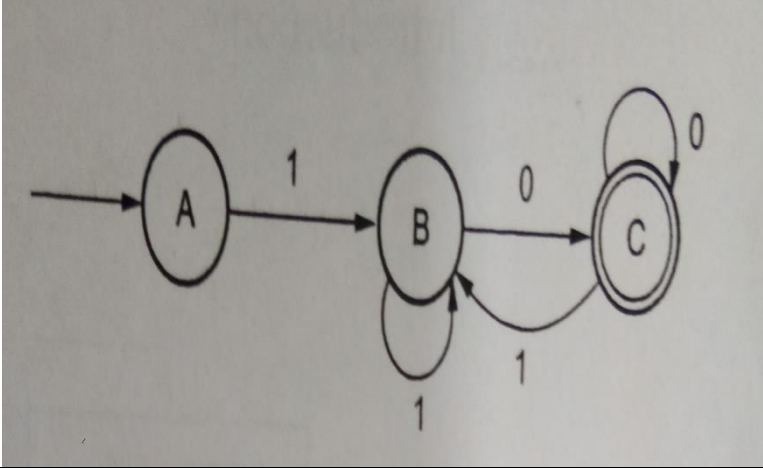
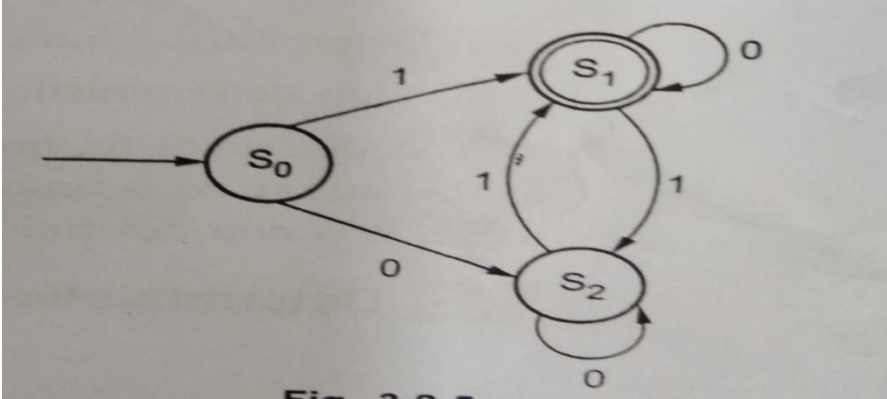
Curriculum Scheme: Rev2016

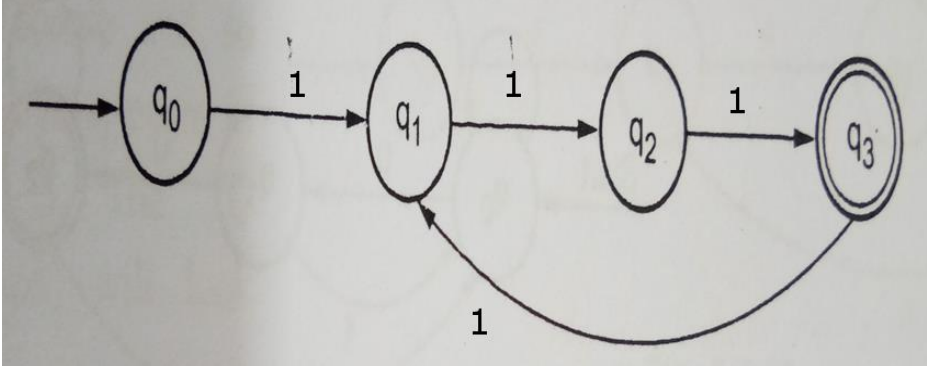
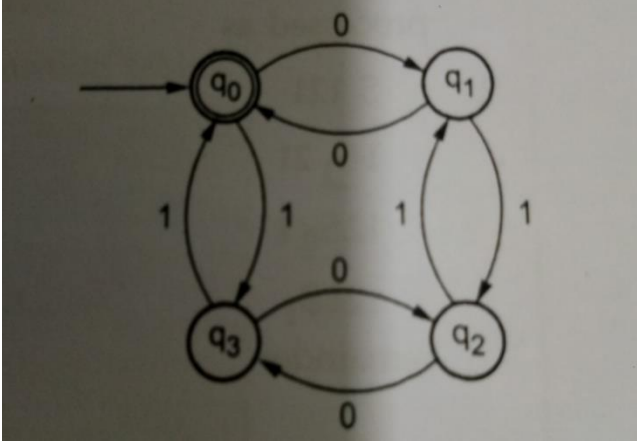
Examination: TE Semester V

Course Code: **CSC504** and Course Name: **Theory of Computer Science**

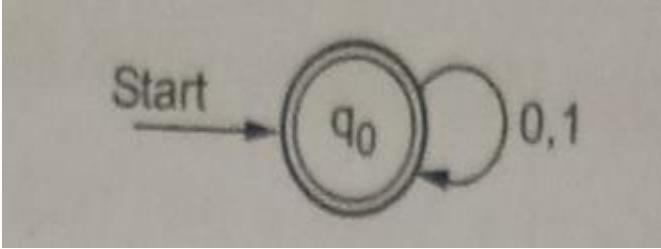
Time: 2 hour

Max. Marks: 80

Q1.	Choose the correct option for following questions. All the Questions are compulsory and carry equal marks
1.	Figure shows finite automata which accepts only those strings _____ 
Option A:	which start with 1 and ends with 0
Option B:	which contains only input 101
Option C:	which start with 1 and ends with 1
Option D:	which start with ϵ and ends with 1
2.	Figure shows finite automata which accepts _____ 

Option A:	odd number of 1's and any number of 0's.
Option B:	odd number of 0's and any number of 1's.
Option C:	even number of 1's and any number of 0's.
Option D:	odd number of 0's and even number of 1's.
3.	<p>Figure shows finite automata which checks _____</p> 
Option A:	whether the given unary number is divisible by 3
Option B:	whether the given unary number is divisible by 2
Option C:	whether the given unary number is divisible by 4
Option D:	whether the given unary number is divisible by 0
4.	<p>Figure shows finite automata which accepts _____</p> 
Option A:	Even number of 0's and odd number of 1's

Option B:	Odd number of 0's and even number of 1's
Option C:	Even number of 0's and even number of 1's
Option D:	Odd number of 0's and odd number of 1's
5.	<p>Following NFA with ϵ represents language consisting_____</p>
Option A:	The strings of any number of a's followed by any number of b's followed by any number of c's
Option B:	The strings of any number of a's followed by any number of ϵ , followed by any number of c's
Option C:	The strings of any number of a's followed by any number of b's followed by any number of ϵ
Option D:	The strings of any number of ϵ followed by any number of b's followed by any number of c's
6.	<p>ϵ-closures of q_0, q_1 and q_2 are obtained as _____ for following NFA with ϵ</p>
Option A:	ϵ -closure(q_0)= $\{q_0\}$, ϵ -closure(q_1)= $\{q_1, q_2\}$, ϵ -closure(q_2)= $\{q_2\}$
Option B:	ϵ -closure(q_0)= $\{q_0, q_1\}$, ϵ -closure(q_1)= $\{q_1, q_2\}$, ϵ -closure(q_2)= $\{q_2\}$

Option C:	$\epsilon\text{-closure}(q_0)=\{q_0, q_1\}$, $\epsilon\text{-closure}(q_1)=\{q_1\}$, $\epsilon\text{-closure}(q_2)=\{q_2\}$
Option D:	$\epsilon\text{-closure}(q_0)=\{q_0\}$, $\epsilon\text{-closure}(q_1)=\{q_1\}$, $\epsilon\text{-closure}(q_2)=\{q_2\}$
7.	<p>Following DFA represents Language _____</p> 
Option A:	Containing any combination of 0 and 1
Option B:	Containing equal number of zeros and 1's
Option C:	Containing all the string except ϵ
Option D:	Containing odd number of 0's and 1's
8.	Regular expression =0(00)* represents the language _____
Option A:	having odd number of 0's
Option B:	having even number of 0's
Option C:	having equal number of 0's
Option D:	having any number of 0's as well as empty string
9.	_____ is the regular expression to denote the language L over the set $\Sigma=\{a,b,c\}$ such that every string will have atleast one a followed by atleast one b followed by atleast one c
Option A:	$a^+ b^+ c^+$
Option B:	$a^* b^* c^*$
Option C:	$a^* b^* c$

Option D:	ab^*c^*
10.	_____ is R.E. for the language L which accepts all the strings with atleast two b's over the set $\Sigma=\{a,b\}$
Option A:	$(a+b)^* b (a+b)^* b (a+b)^*$
Option B:	$(a+b)^* (a+b)^* (a+b)^*$
Option C:	$(a+b)^+ (a+b)^*(a+b)^+$
Option D:	$(a+b) (a+b) (a+b)^*$
11.	Production rules for the CFG for the language having any number of a's over the set $\Sigma=\{a\}$
Option A:	$S \rightarrow aS$ and $S \rightarrow \epsilon$
Option B:	$S \rightarrow aS$
Option C:	$S \rightarrow a$
Option D:	$S \rightarrow S$
12.	The rule for _____ is Non terminal=one terminal.Any number of non-terminals
Option A:	GNF
Option B:	CNF
Option C:	Simplified grammer
Option D:	LBA
13.	In _____ we can remove epsilon production, unit production and useless symbol without changing the meaning.
Option A:	Finite Automata

Option B:	Context free grammer
Option C:	Turing machine
Option D:	Linear bounded automata
14.	The grammar $S \rightarrow (S) \mid SS \mid \epsilon$ is not suitable for predictive parsing because the grammar is
Option A:	Right recursive
Option B:	Left recursive
Option C:	Ambiguous
Option D:	An operator grammar
15.	_____ is the instantaneous description to design PDA for accepting language $L = \{a^n b^{2n} \mid n \geq 1\}$
Option A:	$\delta(q_0, a, Z_0) = (q_0, aaZ_0)$ $\delta(q_0, a, a) = (q_0, aaa)$ $\delta(q_0, b, a) = (q_1, \epsilon)$ $\delta(q_1, b, a) = (q_1, \epsilon)$ $\delta(q_1, \epsilon, Z_0) = (q_2, \epsilon)$
Option B:	$\delta(q_0, a, Z_0) = (q_0, aZ_0)$ $\delta(q_0, a, a) = (q_0, a)$ $\delta(q_0, b, a) = (q_1, ba)$ $\delta(q_1, b, a) = (q_1, ab)$ $\delta(q_1, \epsilon, Z_0) = (q_2, \epsilon)$
Option C:	$\delta(q_0, a, Z_0) = (q_0, a)$ $\delta(q_0, a, a) = (q_0, aa)$ $\delta(q_0, b, a) = (q_1, b)$ $\delta(q_1, b, a) = (q_1, a)$ $\delta(q_1, \epsilon, Z_0) = (q_1, Z_0)$
Option D:	$\delta(q_0, a, Z_0) = (q_0, a)$ $\delta(q_0, a, a) = (q_0, aa)$ $\delta(q_0, b, a) = (q_1, ab)$ $\delta(q_1, b, a) = (q_1, ab)$ $\delta(q_1, \epsilon, Z_0) = (q_1, Z_0)$

16.	$L=0^m1^n0^{m+n}$ can be constructed by using _____
Option A:	DFA
Option B:	NFA
Option C:	PDA
Option D:	Moore
17.	Logic to construct turing machine for the language $L=a^n b^n$ where $n \geq 1$ is _____
Option A:	Convert a by A and then move ahead along the input tape and find out the b convert it to B. Repeat this process for all a's and b's
Option B:	Convert b by B and then move ahead along the input tape and find out the a convert it to A.
Option C:	Convert a by A and then move ahead along the input tape and find out the b convert it to B.
Option D:	Convert all a's by A first and then convert all b's to B.
18.	In the high level languages use of _____ built the modularity in the program development process
Option A:	Subroutines
Option B:	Function
Option C:	stack
Option D:	code
19.	Logic to construct TM for the addition function for the unary number system is _____
Option A:	To simply replace + by 1 and move ahead right for searching end of the string and then we will convert last 1 to Δ .

Option B:	To move ahead right for searching end of the string and then we will convert last 1 to Δ .
Option C:	To simply replace + by 1 and move ahead right for searching end of the string Δ .
Option D:	To move ahead right for searching end of the string.
20.	The undecidability of strings is determined with the help of _____
Option A:	Post correspondence theorem
Option B:	Rice theorem
Option C:	halting
Option D:	pre-correspondence theorem

Q2. (20 Marks Each)	Solve any Four out of Six	5 marks each
A	Design a DFA to accept string of a's and b's ending with 'abb' over I/P $z=\{a,b\}$	
B	Design PDA for the language that accepts the strings with $n_a(w) < n_b(w)$ where $w \in (a+b)^*$	
C	Design a mealy machine to find 2's complement of a given binary number.	
D	Remove the ϵ production from following CFG by preserving meaning of it. $S \rightarrow XYX$ $X \rightarrow 0X \epsilon$ $Y \rightarrow 1Y \epsilon$	
E	Construct Turing Machine for $L = a^n b^n c^n \mid n \geq 1$	
F	Write short note on Rice Theorem	

Q3. (20 Marks Each)	Solve any Two Questions out of Three	10 marks each
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A	<p>Consider the following grammar</p> $S \rightarrow iCtS \mid iCtSeS \mid a$ $C \rightarrow b$ <p>For the string ibtibtaea find the following i) Leftmost derivation ii) Rightmost derivation iii) Parse Tree iv) Check if the above grammar is ambiguous.</p>															
B	<p>Convert the following NFA to DFA. P is the initial state and r and s are the final states.</p> <table border="1" data-bbox="459 533 1433 813"> <thead> <tr> <th>δ</th> <th>0</th> <th>1</th> </tr> </thead> <tbody> <tr> <td>P</td> <td>{p, r}</td> <td>{q}</td> </tr> <tr> <td>q</td> <td>{r, s}</td> <td>{p}</td> </tr> <tr> <td>r*</td> <td>{p, s}</td> <td>{r}</td> </tr> <tr> <td>s*</td> <td>{q, r}</td> <td>{ }</td> </tr> </tbody> </table>	δ	0	1	P	{p, r}	{q}	q	{r, s}	{p}	r*	{p, s}	{r}	s*	{q, r}	{ }
δ	0	1														
P	{p, r}	{q}														
q	{r, s}	{p}														
r*	{p, s}	{r}														
s*	{q, r}	{ }														
C	<p>Construct PDA for the grammar</p> $E \rightarrow E + E \mid E - E \mid (E) \mid id$															